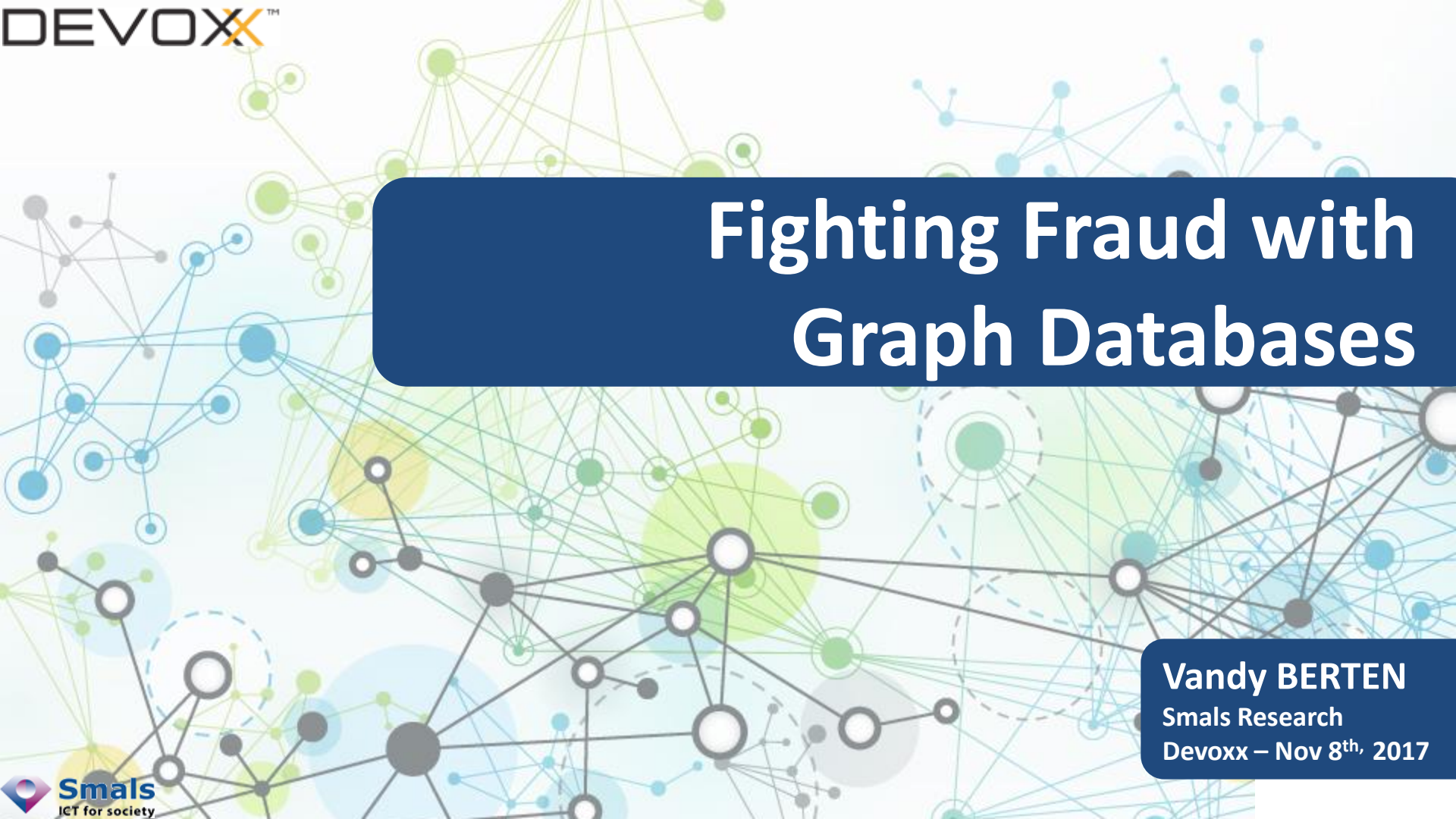
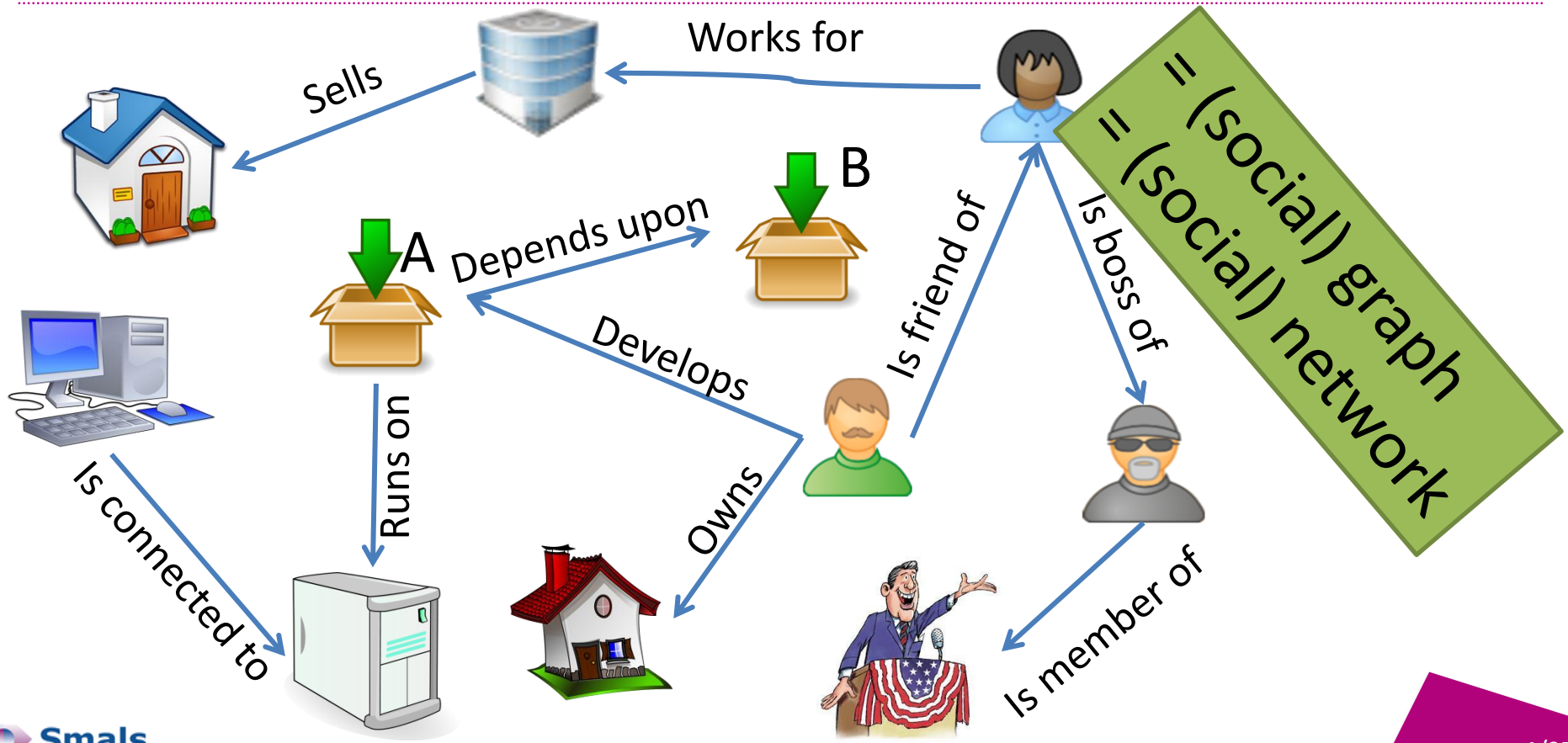

DEVOXX™
Fighting Fraud with Graph Databases
Vandy BERTEN
Smals Research
Devoxx – Nov 8th, 2017
Smals
ICT for society

Outline

MOTIVATIONS

Motivations



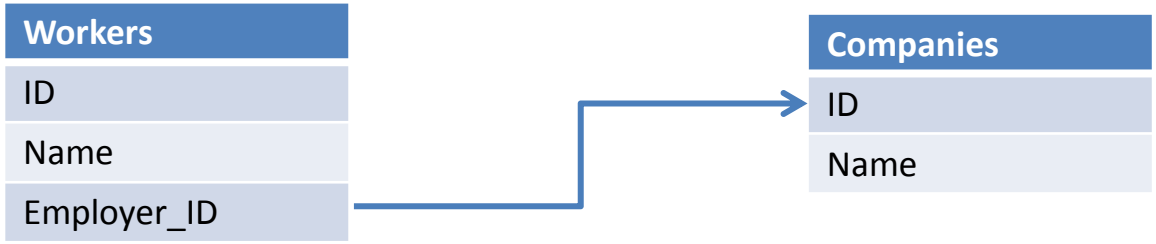
Starting points

- Graphs excel in model fraudulent/criminal behavior
- RDMS perform poorly in representing graphs (i.e., relationships between entities)
- Graph databases propose a new database model, allowing to query graphs

- RBMS focus on entities, Graph DB focus on relationships

GRAPH DATABASES

RDBMS and relationships



Smals' Employees?

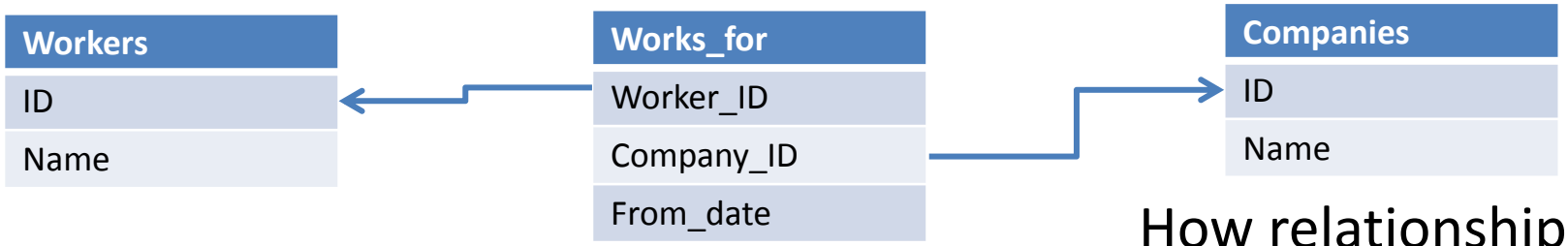
How relationships are implemented

What concerns developers

```
SELECT  Workers.Name
FROM    Workers
JOIN    Companies
  ON    Workers.Employer_ID = Companies.ID
WHERE   Companies.Name = 'Smals'
```

RDBMS : Relational DataBase Management System

RDBMS and relationships



Smals' Employees?

How relationships are implemented

What concerns developers

```
SELECT  Workers.Name
FROM    Workers
JOIN    Works_for
  ON    Workers.ID = Works_for.Worker_ID
JOIN    Companies
  ON    Works_for.Company_ID = Companies.ID
WHERE   Companies.Name = 'Smals'
```

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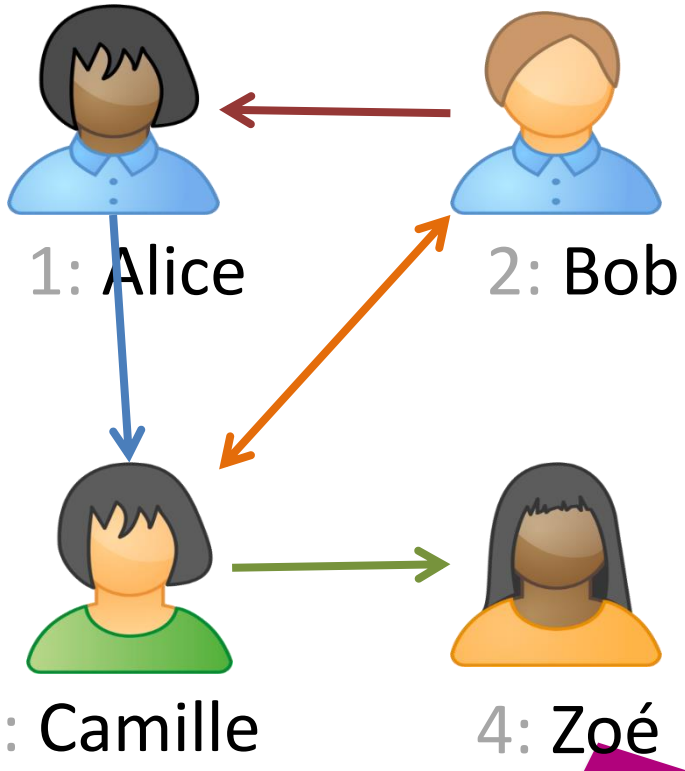
RDBMS and relationships

ID	Name
1	Alice
2	Bob
3	Camille
4	Zoé

Liker_ID	Liked_ID
1	3
2	1
2	3
3	2
3	4




```
SELECT  p1.Name
FROM    People p1
JOIN    Likes
  ON    Likes.Liked_ID = p1.ID
JOIN    People p2
  ON    Likes.Liker_ID = p2.ID
WHERE   p2.Name = "Bob"
```

Who Bob likes?

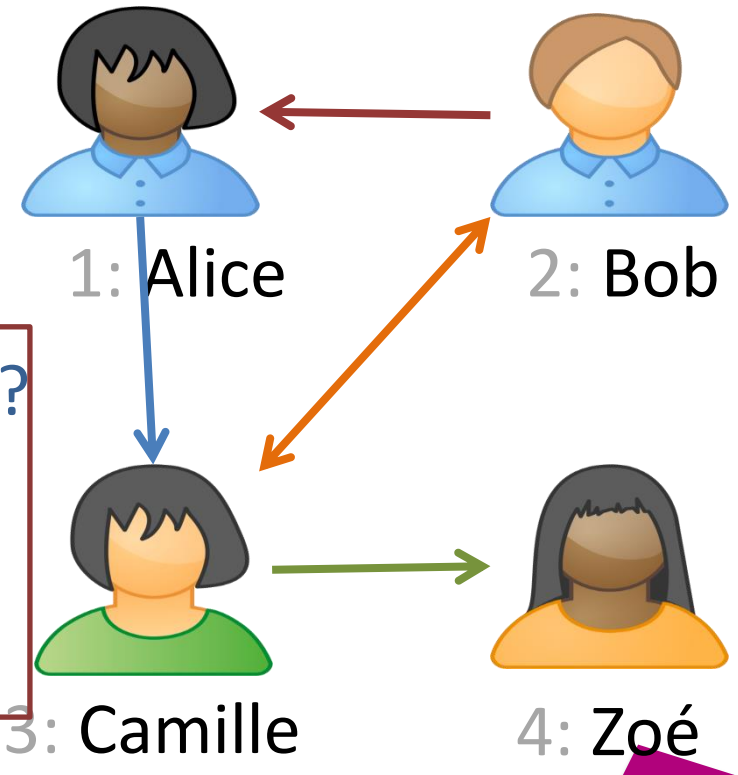
RDBMS and relationships

ID	Name
1	Alice
2	Bob
3	Camille
4	Zoé

Liker_ID	Liked_ID
1	3
2	1
2	3
3	2
3	4

```
SELECT  p1.Name
FROM    People p1
JOIN    Likes l1
   ON   l1.Liked_ID = p1.ID
JOIN    Likes l2
   ON   l1.Liker_ID = l2.Liked_ID
JOIN    People p2
   ON   l2.Liker_ID = p2.ID
WHERE   p2.Name = "Bob"
```

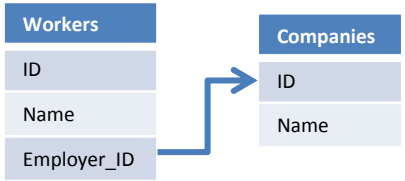
Who Bob (likes)² ?



RDBMS and relationships: limitations

1-n relationships

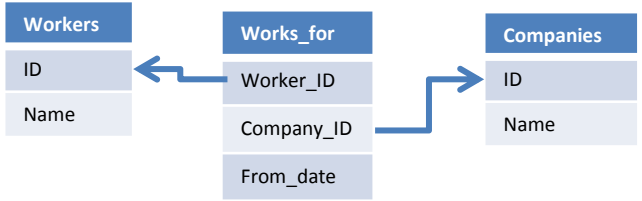
Foreign key:



- Transforms the role of attributes
- No type « Key »

m-n relationships

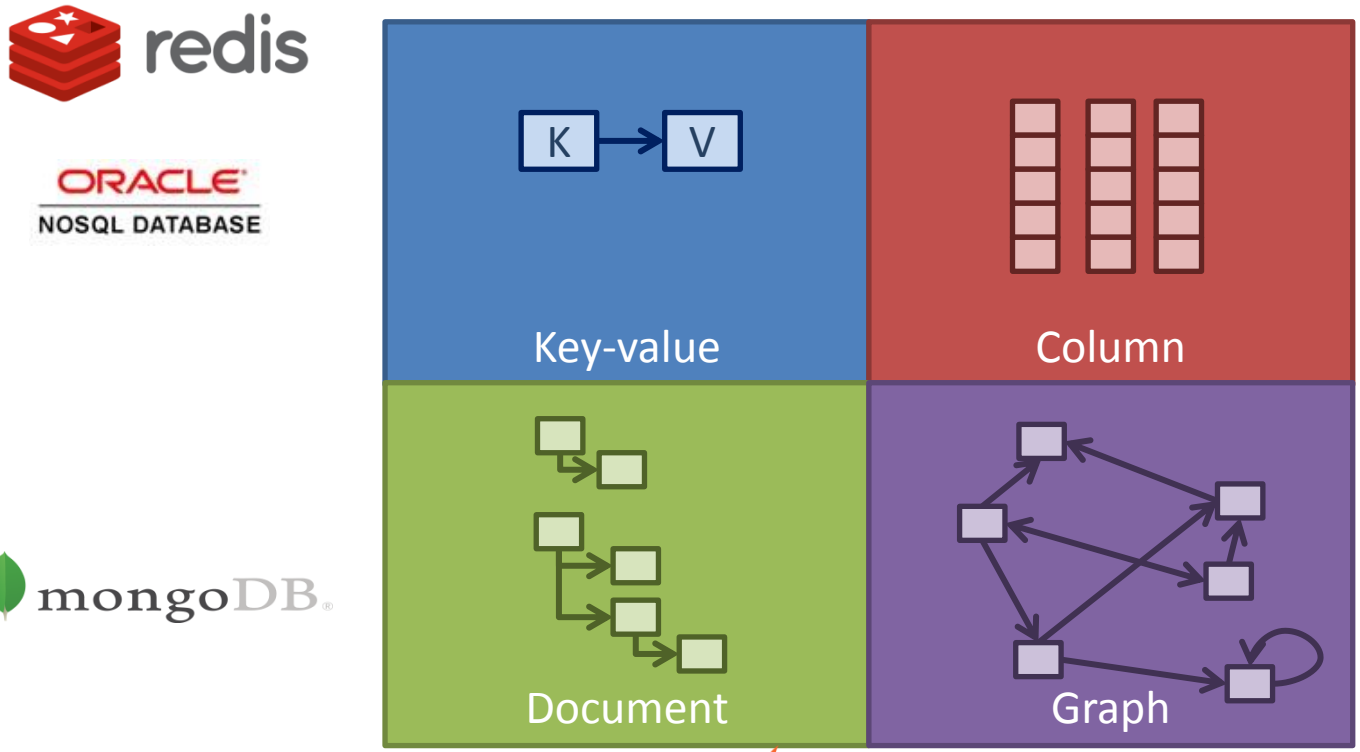
Join table:



- Transforms the role of tables
- No distinguishable

- Integrity constraints possible (extern to tables)
- Relationship structure to be given to each request
- Join : Heavy in writing code and in execution time

NoSQL solutions



redis

ORACLE NOSQL DATABASE

cassandra

Google BigTable

APACHE HBASE

AllegroGraph

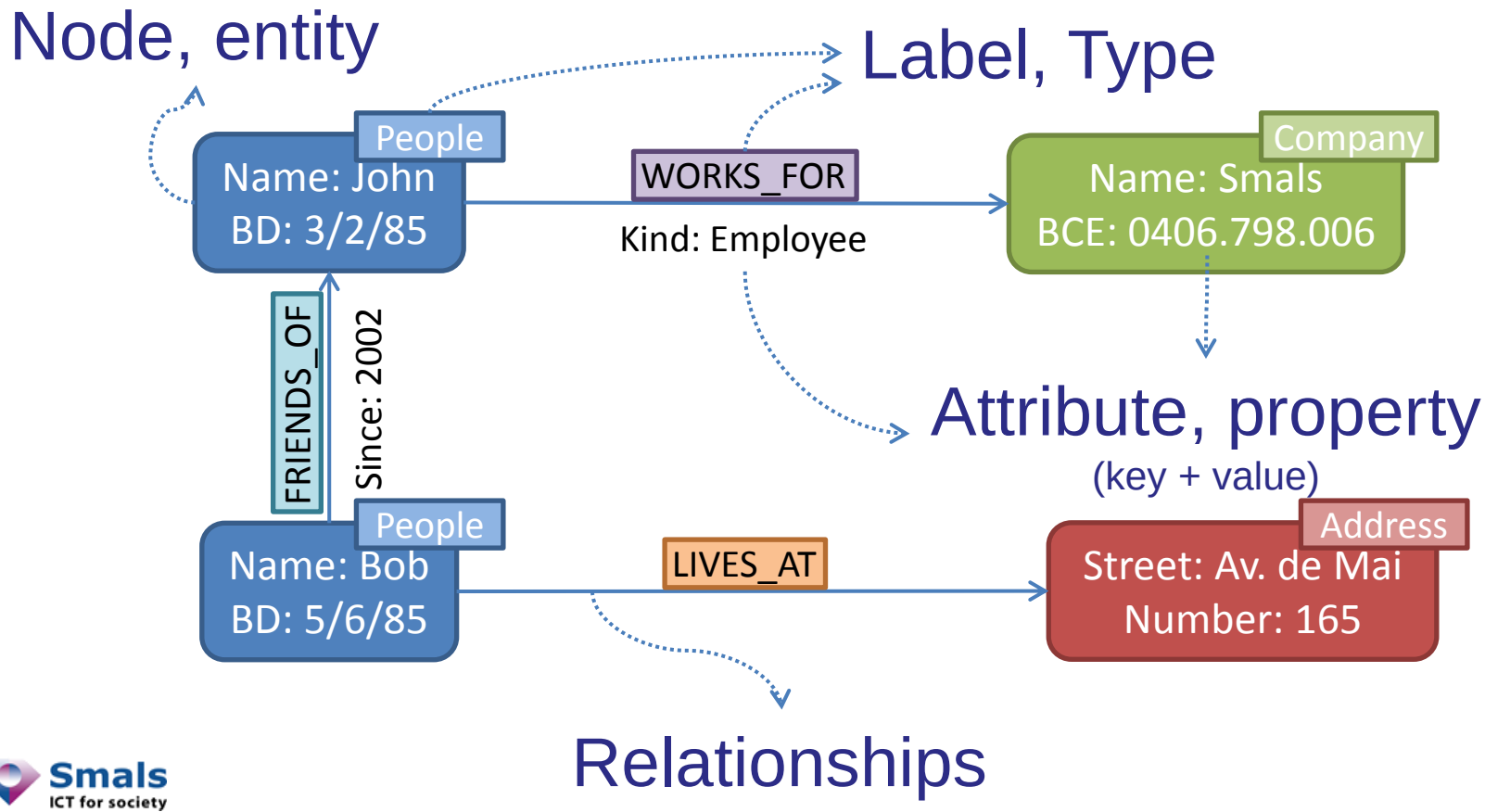
mongoDB

neo4j

InfiniteGraph
Powered by Objectivity

OrientDB

Graph database model



GraphDB: objectives

Objectives of graph databases:

Querying language splitting apart relationship definition (creation time) and « **matching** »

Fast **search engine** for **relationship traversals**

```
SELECT p1.Name
FROM   People p1
JOIN   Likes l1
   ON  l1.Liked_ID = p1.ID
JOIN   Likes l2
   ON  l1.Liker_ID = l2.Liked_ID
JOIN   People p2
   ON  l2.Liker_ID = p2.ID
WHERE  p2.Name = "Bob"
```

Cypher (Neo4j)

```
MATCH
   (:People {Name:"Bob"})
   -[:Likes*2]->
   (p:People)
RETURN p.Name
```

GraphDB: less code

```
(SELECT T.directReportees AS directReportees, sum(T.count) AS count
FROM (
SELECT manager.pid AS directReportees, 0 AS count
  FROM person_reportee manager
  WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
UNION
  SELECT manager.pid AS directReportees, count(manager.directly_manages) AS count
FROM person_reportee manager
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
GROUP BY directReportees
UNION
SELECT manager.pid AS directReportees, count(reportee.directly_manages) AS count
FROM person_reportee manager
JOIN person_reportee reportee
ON manager.directly_manages = reportee.pid
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
GROUP BY directReportees
UNION
SELECT manager.pid AS directReportees, count(L2Reportees.directly_manages) AS count
FROM person_reportee manager
JOIN person_reportee L1Reportees
ON manager.directly_manages = L1Reportees.pid
JOIN person_reportee L2Reportees
ON L1Reportees.directly_manages = L2Reportees.pid
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
GROUP BY directReportees
) AS T
GROUP BY directReportees)
UNION
(SELECT T.directReportees AS directReportees, sum(T.count) AS count
FROM (
SELECT manager.directly_manages AS directReportees, 0 AS count
FROM person_reportee manager
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
UNION
SELECT reportee.pid AS directReportees, count(reportee.directly_manages) AS count
FROM person_reportee manager
JOIN person_reportee reportee
ON manager.directly_manages = reportee.pid
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
GROUP BY directReportees
UNION
```

```
SELECT depth1Reportees.pid AS directReportees,
count(depth2Reportees.directly_manages) AS count
FROM person_reportee manager
JOIN person_reportee L1Reportees
ON manager.directly_manages = L1Reportees.pid
JOIN person_reportee L2Reportees
ON L1Reportees.directly_manages = L2Reportees.pid
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
GROUP BY directReportees
) AS T
GROUP BY directReportees)
UNION
(SELECT T.directReportees AS directReportees, sum(T.count) AS count
  FROM(
  SELECT reportee.directly_manages AS directReportees, 0 AS count
FROM person_reportee manager
JOIN person_reportee reportee
ON manager.directly_manages = reportee.pid
WHERE manager.pid = (SELECT id FROM person WHERE name = "fName lName")
GROUP BY directReportees
UNION
SELECT L2Reportees.pid AS directReportees, count(L2Reportees.directly_manages)
AS count
FROM person_reportee manager
JOIN person_reportee L1Reportees
```

```
MATCH
   (boss)-[:MANAGES*0..3]->(sub),
   (sub)-[:MANAGES*1..3]->(report)
WHERE boss.name = « John Doe"
RETURN sub.name AS Subordinate,
       count(report) AS Total
```

GraphDB : less code

"We found Neo4j to be literally thousands of times faster than our prior MySQL solution, with queries that require 10-100 times less code. Today, Neo4j provides eBay with functionality that was previously impossible."

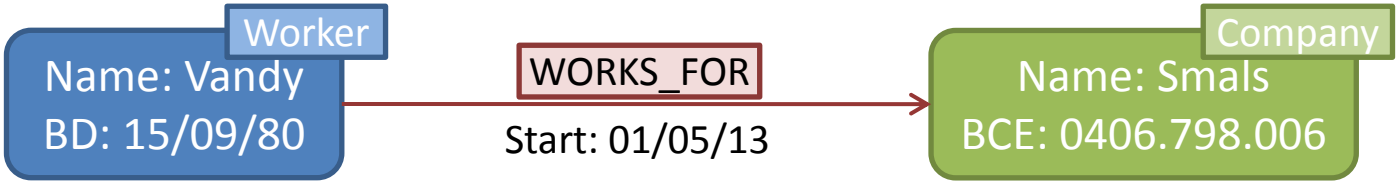
https://neo4j.com/news/graph-databases-provide-crystal-ball-online-recommendations-wal-mart-ebay/

NEO4J

Neo4j

- Current Graph Database leader
- Written in Java, multi-plateform
- Community Edition (free, GPL) + Enterprise Edition ($, AGPL)
- Can be interfaced by many languages
- ACID-compliant
- Queries in Cypher or Gremlin
- Since 2007

Neo4j request



Node of type « Worker »

Node of type « Company »

Relationship of type « WORKS_FOR »

```
MATCH   (w:Worker)   -[r:WORKS_FOR]->   (c:Company)
WHERE   c.Name = "Smals"
RETURN  w.Name, w.BD, r.Start
```

Variante :

```
MATCH
  (w:Worker)
  -[r:WORKS_FOR]->
  (c:Company {Name:"Smals"})
RETURN …
```

Smals' employees

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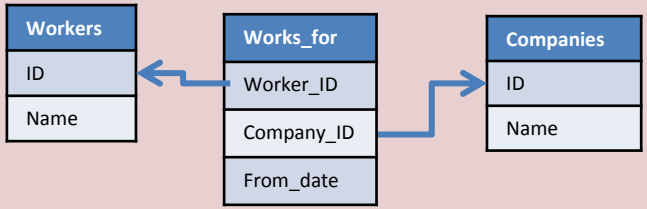
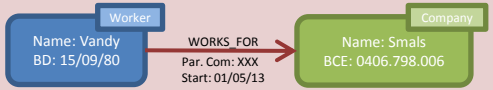
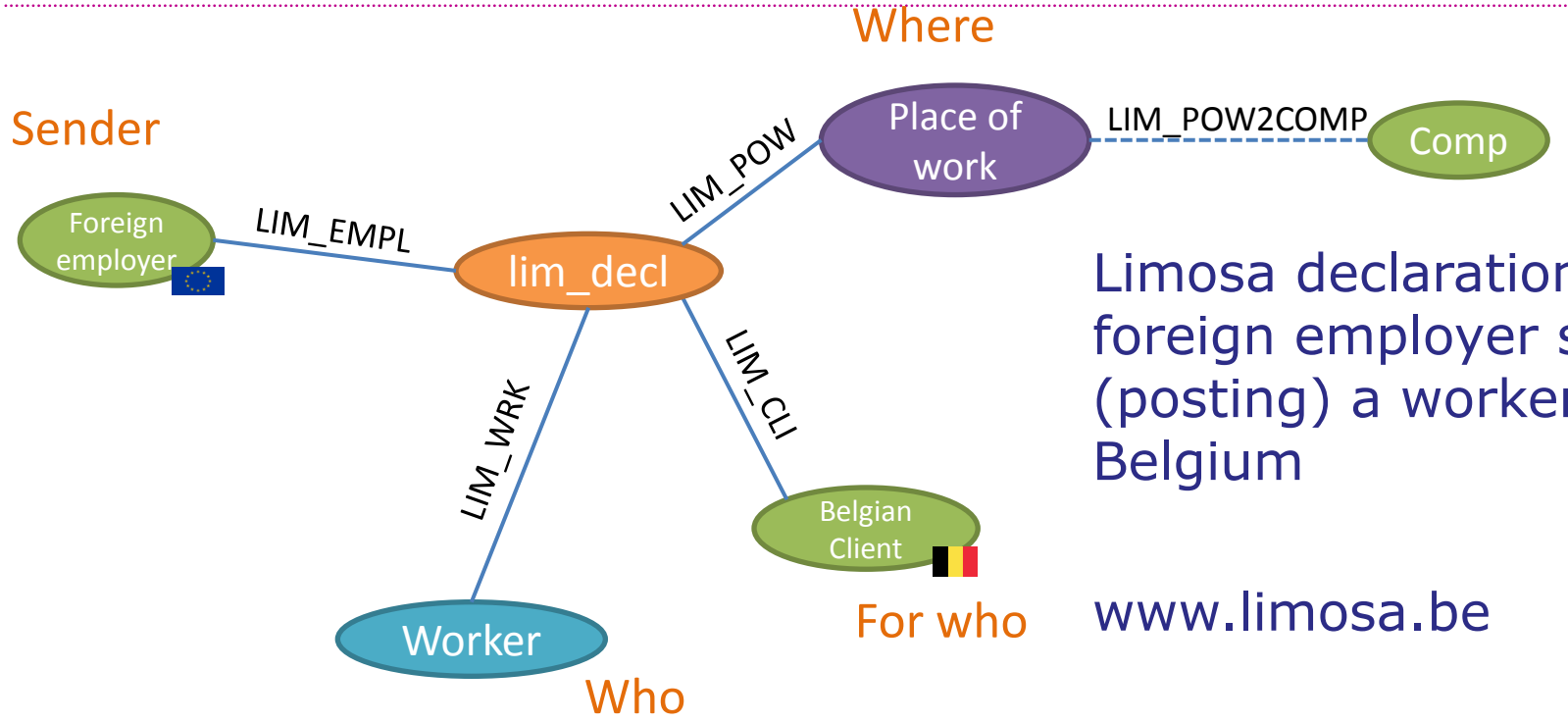
How it works

	RDBMS	Graph DB
schema		
WHERE …	Search in Companies (with index): O(log([Table size])	Search in nodes (with index) : O(log([DB size])
Relation	JOIN : - Search in Works_for : O(log([Table size])) - Search in Workers : O(log([Table size]) x [nb res])	Follow pointers in node : O([nb of employees])

USE CASE: SOCIAL DUMPING

Limosa: detachment declaration



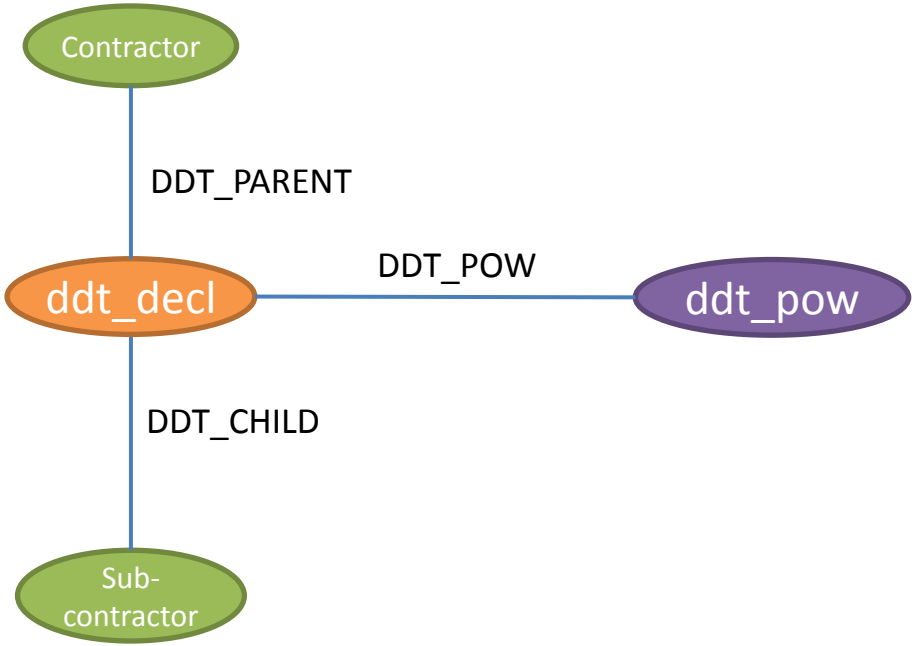
Limosa declaration: A foreign employer sending (posting) a worker in Belgium

www.limosa.be

Limosa : Landenoverschrijdend Informatiesysteem ten behoeve van MigratieOnderzoek bij de Sociale Administratie

Syst. d'infor. transfrontalier en vue de la recherche en matière de migration auprès de l'adm. sociale

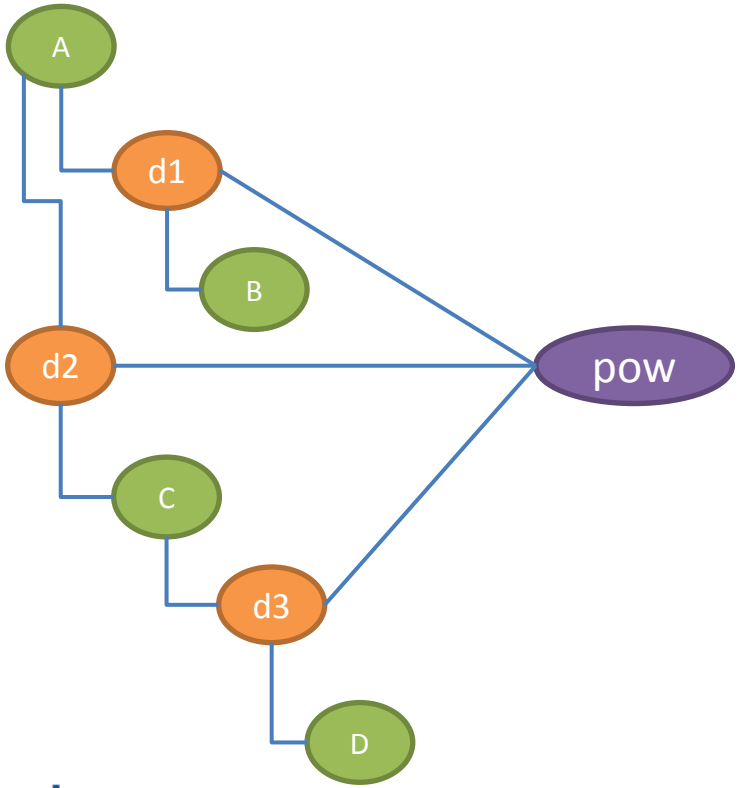
DDT: sub-contractors



DDT: Déclaration De Travaux

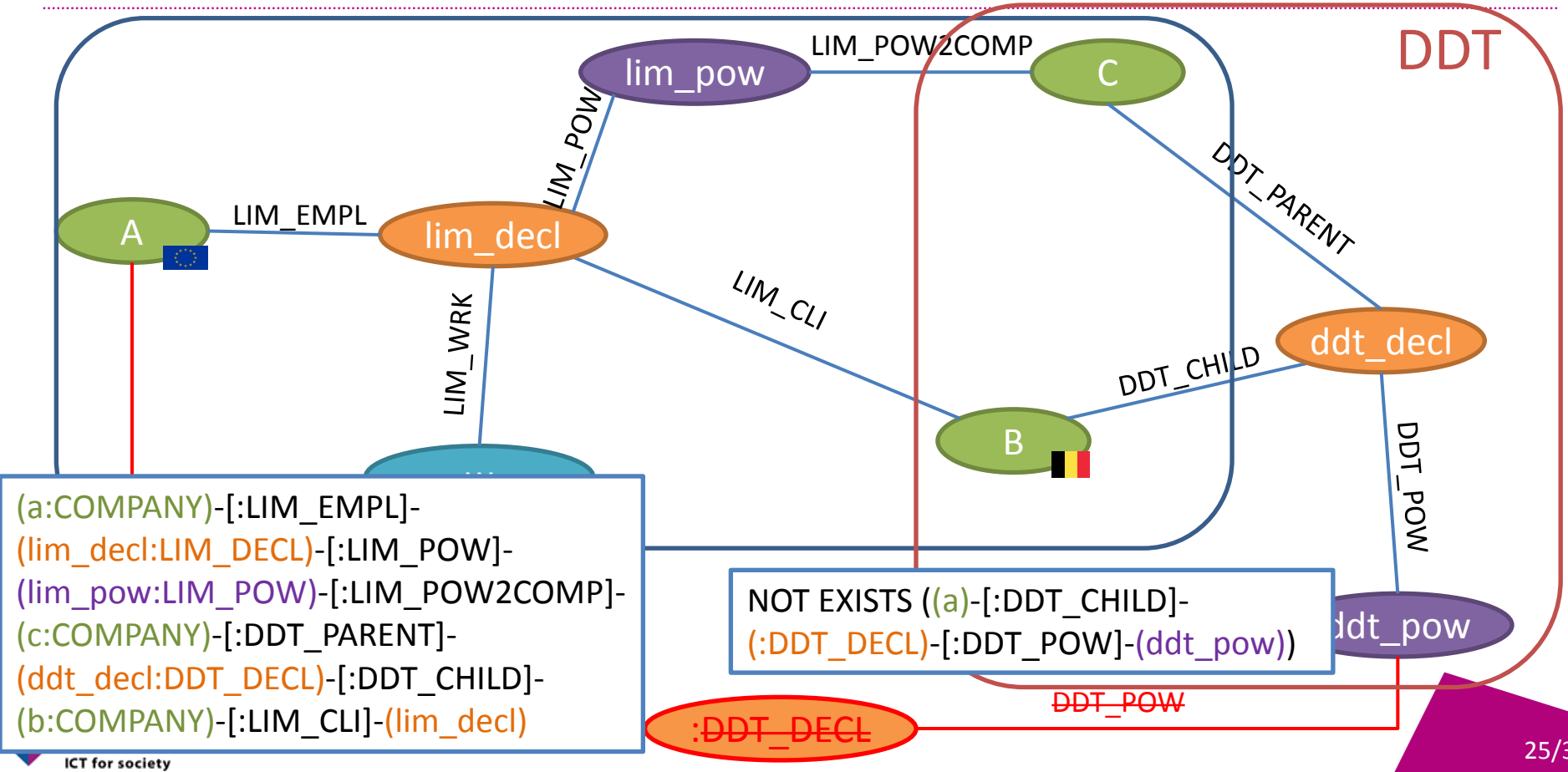
Aangifte van werken

DDT: sub-contractors



On a specific worksite (PoW), we have a chain/tree of sub-contractors

Full schema


DDT
lim_pow
LIM_POW2COMP
C
LIM_POW
DDT_PARENT
A
LIM_EMPL
lim_decl
LIM_CLI
LIM_WRK
ddt_decl
DDT_CHILD
B
DDT_POW
(a:COMPANY)-[:LIM_EMPL]-
(lim_decl:LIM_DECL)-[:LIM_POW]-
(lim_pow:LIM_POW)-[:LIM_POW2COMP]-
(c:COMPANY)-[:DDT_PARENT]-
(ddt_decl:DDT_DECL)-[:DDT_CHILD]-
(b:COMPANY)-[:LIM_CLI]-(lim_decl)
NOT EXISTS ((a)-[:DDT_CHILD]-
(:DDT_DECL)-[:DDT_POW]-(ddt_pow))
ddt_pow
DDT_POW
:DDT_DECL

Full query

```
MATCH
 (a:COMPANY)-[:LIM_EMPL]-(lim_decl:LIM_DECL)-[:LIM_POW]-
 (lim_pow:LIM_POW)-[:LIM_POW2COMP]-(c:COMPANY)-[:DDT_PARENT]-
 (ddt_decl:DDT_DECL)-[:DDT_CHILD]-(b:COMPANY)-[:LIM_CLI]-(lim_decl),
 (ddt_decl)-[:DDT_POW]-(ddt_pow:DDT_POW)
WHERE
 NOT EXISTS ((a)-[:DDT_CHILD]-(:DDT_DECL)-[:DDT_POW]-(ddt_pow)) …
RETURN
 DISTINCT a, b, c,
 COUNT(DISTINCT lim_decl),
 COUNT(DISTINCT ddt_decl)
ORDER BY COUNT(DISTINCT lim_decl) DESC
```

Implementation

- We copied in Neo4j 17M nodes, 45M relationships :
  - 800.000 companies, 1.650.000 (foreign) workers
  - 13 millions declarations
  - 185.000 duplication groups
  - 250.000 streets
  - …
- 9 node labels, 15 relationship types
- All companies triplets matching 1st pattern: ~10 seconds (325 results)

In SQL/RDBMS ?

- In SQL/RDBMS,
  - Any node and any relationship is a record : ~65 M records
  - Any type/label corresponds to a table : 24 tables
- Any (n:Type1)-[:REL1]-(m:Type2) corresponds to :

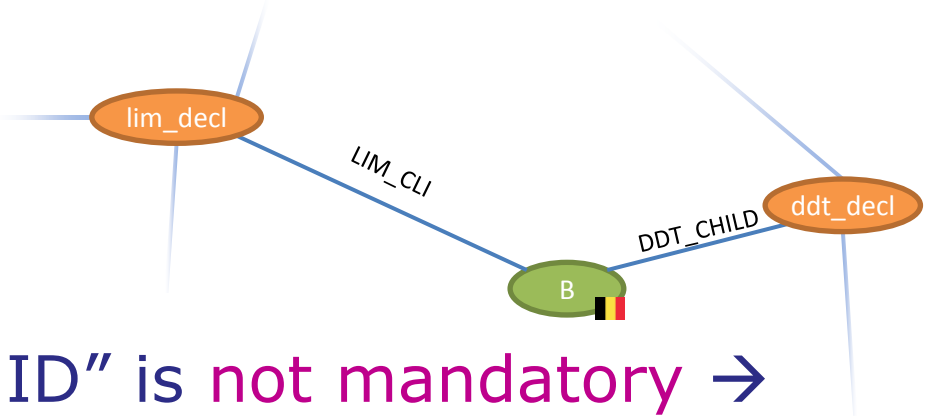
```
FROM Type1
JOIN REL1  ON Type1.key = REL1.type1_key
JOIN Type2 ON Type2.key = REL1.type2_key
```

→ Almost 20 joins!

- → Query far more complex to write
- → Much longer to run

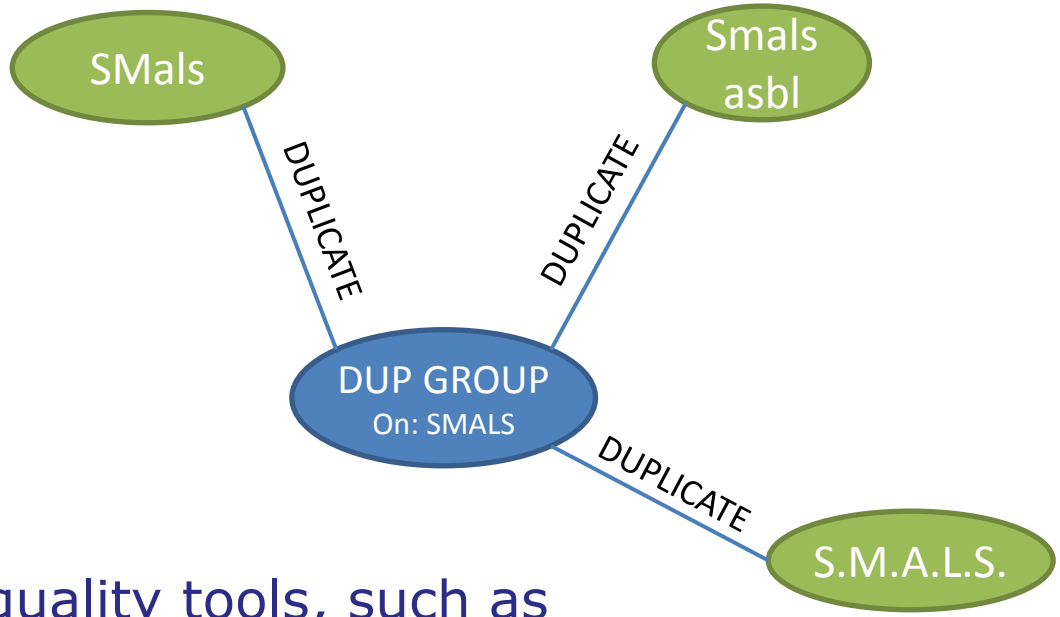
Data quality problem

- This only works if « B » has been declared with the same ID in both systems !



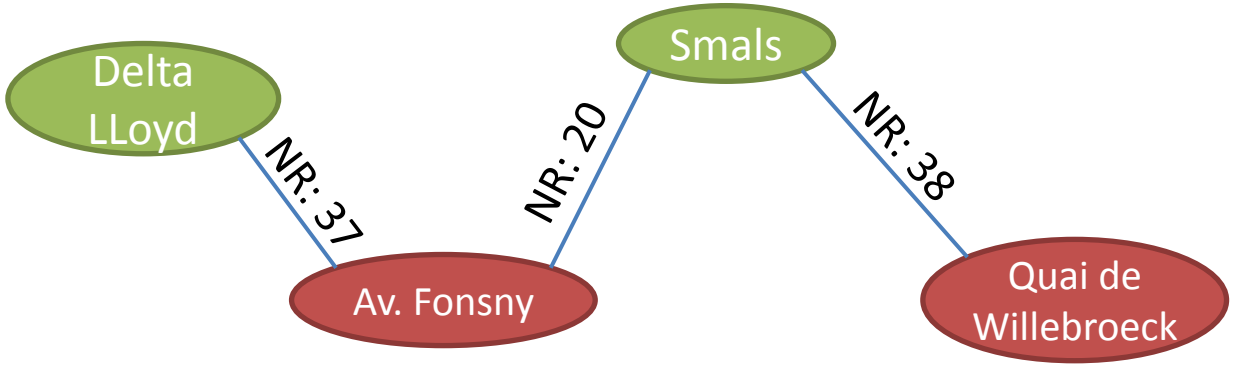
- Sometimes, the “official ID” is not mandatory → name + address is OK
- We have to deal with this!

Duplicates



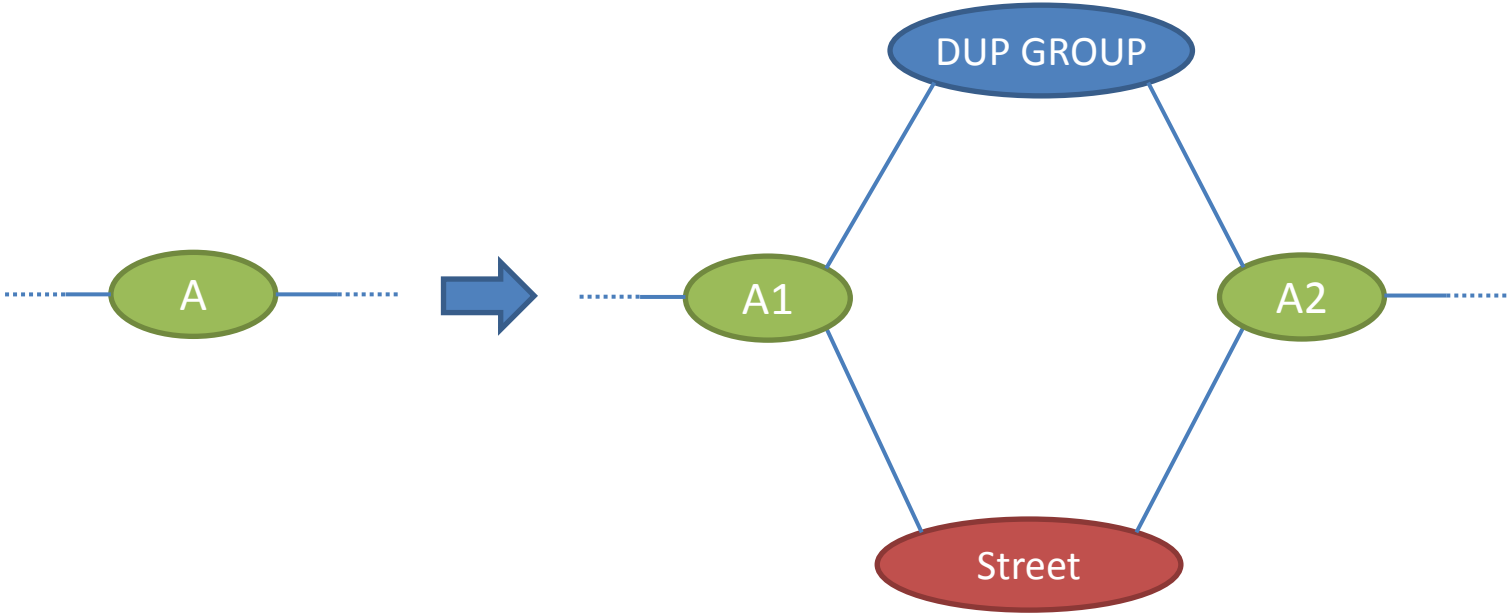
Thanks to data quality tools, such as Trillium (IntoDQ), OpenRefine…

Addresses

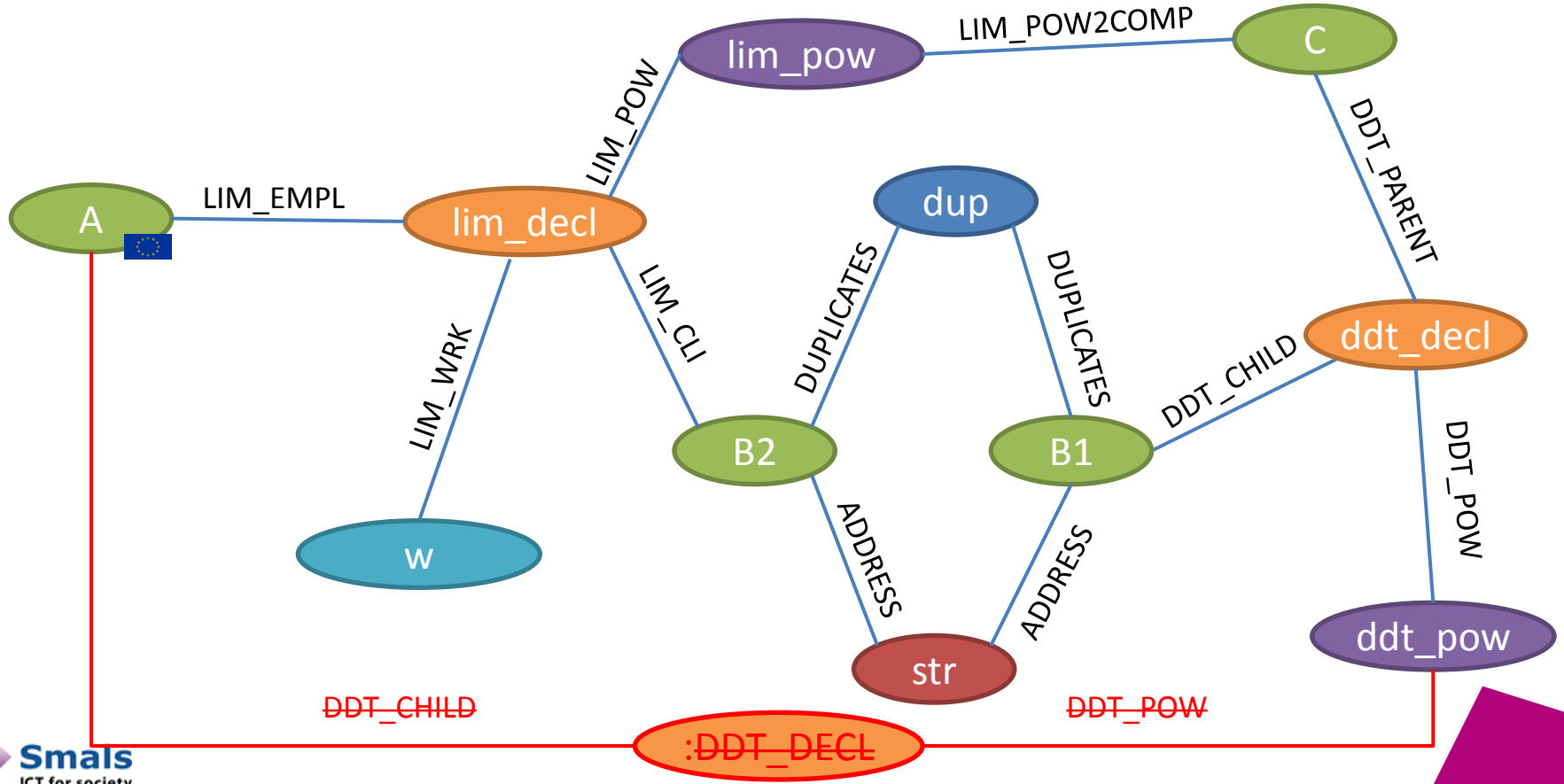


A standardization is needed! Thanks to data quality tools

Duplicates combinations



Schema with B duplicated



Conclusions

- Graph DBs are the « ideal partner » for many fraud detection problems
- They outperform RDBMS on many aspects (query expressiveness, execution speed…)
- Data quality must be taken into consideration
- Not only for fraud: Infrastructure management, recommendation systems, intelligence…

References : www.smalsresearch.be

- Blogs (in French):
  - https://www.smalsresearch.be/un-fraudeur-ne-fraude-jamais-seul/
  - https://www.smalsresearch.be/un-fraudeur-ne-fraude-jamais-seul-partie-2/
  - https://www.smalsresearch.be/bases-de-donnees-relationnelles-adequates-pour-des-relations/
  - https://www.smalsresearch.be/graph-db-vs-rdbms/
- Other:
  - NoSQL - Hype ou Innovation, Grégory Ogonowski, https://www.smalsresearch.be/publications/document/?docid=59

References

- **Graph Databases**, *Robinson & all*, O'Reilly 2015
- **Fraud Detection: Discovering Connections with Graph Databases**, *Neo4j*, Sodawksi & Rathle
- **Fraud Analytics Using Descriptive, Predictive, and Social Network Techniques: A Guide to Data Science for Fraud Detection**, *Bart Baesens, Veronique Van Vlasselaer, Wouter Verbeke*, Willey, 2015

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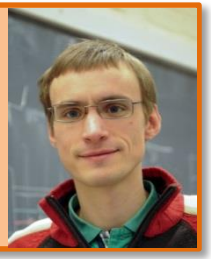

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